

Development of Flowshop Scheduling Simulator

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Abstract

Scheduling or production scheduling is an important activity in production process. Production scheduling is an activity that determine the efficiency of production process. With a good production scheduling, the production process will be running smooth and benefit in cost. With the growing scale in number of jobs and machines on production floor, manual scheduling is irrelevant these days. So a software is required to do a scheduling activity and simulating the result. Scheduling is an activity of allocating resources or machines to run some tasks in specific time range[1]. The result of scheduling is a sequence of jobs that used for reference in determining on how to run a production process. There are several criteria in scheduling process, one of them is minimizing make span [2] or the total amount of time to process all the jobs. The production process has two major flow type, they are flowshop and jobshop.

Flowshop production flow is a production flow type which a sequence of jobs processed on an array of machine. In flowshop production flow every job on the sequence will be processed on every machine on the machine's array, while on jobshop, every job is processed based on the requirement and characteristic of the job. Scheduling in flowshop production flow can be done using many methods. The flow shop can be done using Nawaz Ensore Ham (NEH)[3] method, Palmer[4] method, and Gupta[5] method. Those method prioritizing a job with biggest total of processing time to be processed first.

The previous paper [6] compare the performance of these three methods. In this paper, we develop the flowshop scheduling simulator to make it easy to visualize the running process of these algorithm. The software simulates the result of flowshop scheduling process in a form of gantt chart animation. Gantt chart is used for simulating the flow of production scheduling process. By using simulation, user can see the whole scheduling planning, estimate the completion time of each job, see the total processing time or makespan, and know the tardiness of each job. Similar with the previous paper [6], NEH algorithm constantly works best compare with Palmer and Gupta algorithm in various types of problems.

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I. INTRODUCTION

In textile industries, n jobs $J = \{J_1, J_2, J_3, \dots, J_n\}$ come into queues to be processed in m machines $M = \{M_1, M_2, M_3, \dots, M_m\}$. Each machine has its own functionality. For example cutting, sewing, folding, and

packing machines. Each job has to be processed in each of these machines in a certain sequence. For example, first each job has to go to the cutting machine. After it is done, the job must go to be processed next in the sewing machine, and then to the folding machine, and finally to the packing machines. Each of these jobs has its unique identifier, arrival time, due

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date, and processing time in each machine. The goal is how to arrange all of these jobs in the queues so that it can optimize some objectives. The objectives can be to minimize the total processing time of all jobs in all machines (it's called makespan) and the total tardiness of all jobs.

Scheduling is the art of planning so that the objectives can be achieved. In this research, the objective is to minimize the total processing time of all jobs to be done (makespan) and to minimize the tardiness of all jobs. If the makespan can be minimized, it can reduce the total running cost of a textile industry. If the tardiness of each job can be reduced, it can increase the satisfaction of the customers.

The rule to create the schedule is called heuristics. The previous paper [6] compare some heuristics used in textile industries. In this paper, three heuristics are used as the based to create the flowshop process simulation, namely: NEH, Palmer, and Gupta algorithm.

II. FLOWSHOP SCHEDULING PROBLEM

Flowshop scheduling is a class of scheduling in which there are n jobs to be processed in m machines. Each job must be processed in each of these machine based on a certain sequence. Flowshop problems are characterized by the processing time PT_{ij} of jobs i on machine j , where $1 \leq i \leq n$ and $1 \leq j \leq m$. Each machine m_j can only process one job j_i at a time. Each job can only be processed by one machine at a time. Once a machine has started processing a job, it will continue running on that job until the job is finished.

Given:

- n jobs $J_1, J_2, J_3, \dots, J_n$
- n respective due date $d_1, d_2, d_3, \dots, d_n$

Goal:

- Minimize makespan
- Minimize tardiness

Fig. 1. shows the illustration of flowshop scheduling problem with four types of machines, and 5 jobs that must be done in each of these machines. The processing sequence is start from the machine 1, continue to machine 2, machine 3, and finalized in machine 4. The value along the X axis shows the timeline of the scheduling process. Each job is represented as a single color in the chart and must be processed sequentially from machine 1 until machine 4 to be complete.

For example, job 1 start at time 0 in machine 1, and continue to machine 2 at time 26, and then start at time 60 machine 3, and continue in machine 4 until it's complete at time 88. Job 2 can only be processed in machine 1, after job 1 is finished. The same rule apply to all the jobs in the queue. Based on this constraint, the total time needed to finish all the job in 4 machines is at time 522. The total time to process all the jobs in queue is called makespan.

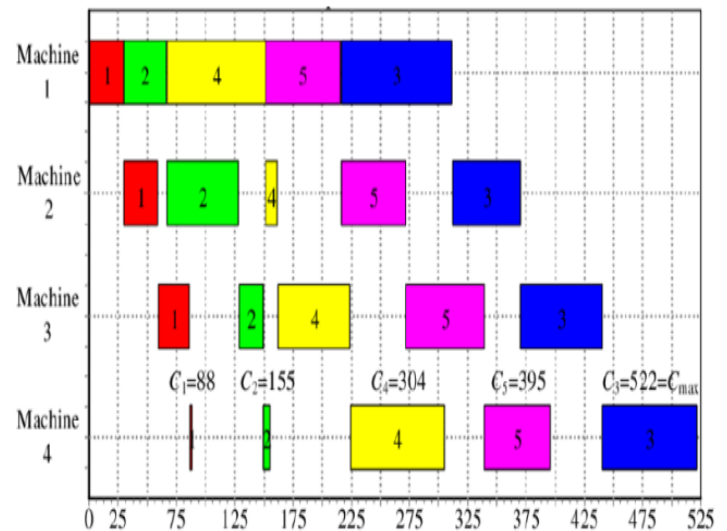


Fig. 1. Flowshop scheduling with 4 machines and 5 jobs

III. FLOWSHOP SCHEDULING ALGORITHM

There are various algorithm to solve flowshop scheduling problem. Three algorithm discussed in the previous paper [6] are: NEH (Nawaz Ensore Ham), Palmer, and Gupta.

III.1 Nawaz Ensore Ham (NEH) Algorithm

This algorithm is proposed by Muhammad Nawaz, E. Emory Ensore Jr, Inyong Ham in 1983 [3] and has awarded as the best algorithm in permutation flowshop problem. This algorithm is based on the longest processing time heuristics. It means it choose the job with the longer processing time to be executed first than the other job with shorter processing time. The complete algorithm can be seen in Algorithm 1.

Algorithm 1: NEH Algorithm

1. Add the processing time of a job on machine 1 until machine m . Then it will be n values show the total processing time of n jobs.
2. Sort decending this n total processing time of n jobs.
3. Start to process the job with the longest total processing time to be executed first.
 - a. Enumerate all the possibilities permutation
 - b. Calculate the sparsial makespan for all the permutation possibilities
 - c. Choose the arrangement with the minimum makespan.

III.2 Palmer Algorithm

This algorithm is proposed by Palmer in 1965, and works based on a value called slope index. This value determine which job should be executed first.

Algorithm 2: Palmer Algorithm

1. Calculate the slope index of each job with this formula:

$$SI_i = \sum_{k=1}^m (m - (2 \times k - 1)) \times PT_{ik}$$

where

SI_i : slope index of job J_i

m : total number of machine

k : step number k of job J_i

PT_{ik} : processing time of job J_i step k

2. Sort descending the jobs in the queue start from the largest to the lowest slope index.
3. Process each job based on the sequence produce in step 2.

III.3 Gupta Algorithm

This algorithm is proposed by Gupta in 1972. Similar like Palmer, Gupta algorithm works based on slope index, except it is weighted.

Algorithm 3: Gupta Algorithm

1. Calculate the weight of each job (e_i) based on this formula:

-if $PT_{i1} \geq PT_{im}$ then $e_i = -1$
-if $PT_{i1} < PT_{im}$ then $e_i = 1$

2. Calculate the slope index of each job with this formula:

$$SI_i = \frac{e_i}{\min(PT_{i1} + PT_{i2}, \dots, PT_{i(m-1)} + PT_{im})}$$

where

SI_i : slope index of job J_i

m : total number of machine

PT_{im} : processing time of job J_i step m

e_i : weight of job J_i

3. Sort descending the jobs in the queue start from the largest to the lowest slope index.
4. Process each job based on the sequence produce in step 2.

IV. THE OBJECTIVES

The objective of the scheduling in this research is to minimize tardiness and makespan. The formula to calculate tardiness is:

$$T_i = \max\{C_i - d_i, 0\} \quad (1)$$

where

C_i : the completion time or the earliest time at which job J_i is completely processed.

d_i : due date of job J_i

$C_i - d_i$: the lateness of job J_i

So tardiness is the maximum lateness of all jobs in the queue.

Makespan is maximum completion time of all jobs in the queue. The formula to calculate makespan is:

$$C_{max} = \max\{C_1, C_2, C_3, \dots, C_n\} \quad (2)$$

where

C_i : the completion time or the earliest time at which job J_i is completely processed.

V. RESULTS AND DISCUSSION

This paper use Taillard's benchmark problem [7] as the study cases, i.e. problems with 5, 10, and 20 machines and from 20 to 500 jobs. Taillard's benchmark problem size is greater than that of the rare examples published. Such sizes correspond to real dimensions of industrial problems.

Some experiments are conducted to see the result of execution process of NEH, Palmer, and Gupta algorithms. These experiments use Taillard's benchmark flowshop scheduling problems [7] that consist of various numbers of jobs and machines.

V.1. Flowshop Scheduling Simulator

Fig. 2 until Fig. 9 show the result of flowshop scheduling simulator developed by using NEH, Palmer, and Gupta algorithms. The result is in the form of animation that shows the simulation of the whole process from time 0 until all the jobs are finished to be processed.

V.2. Flowshop Scheduling Statistics

Fig. 2 until Fig. 4 show the simulation result of flowshop scheduling developed by NEH, Palmer, and Gupta algorithms with 10 machines and 50 jobs. Fig. 5 summarize Fig. 2 until Fig. 4 in the report form.

Fig. 6 until Fig. 8 show the simulation result of flowshop scheduling developed by NEH, Palmer, and Gupta algorithms with 20 machines and 100 jobs. Fig. 9 summarize Fig. 6 until Fig. 8 in the report form.

Only two types of problems are shown in this paper. In fact, several experiments with various number of jobs and machines are conducted to see these three algorithms' performances. Similar with the previous paper [6], experiments show that NEH algorithm perform constantly the best, followed by Palmer and Gupta algorithm.

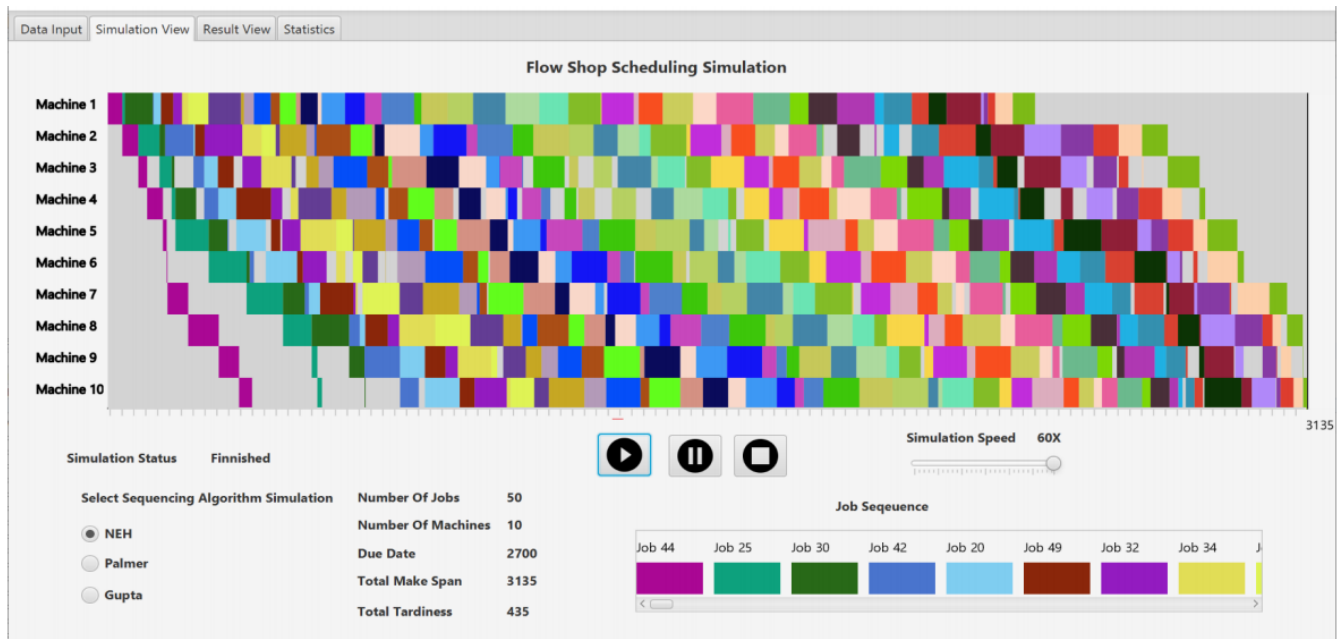


Fig. 2. Final result of NEH execution process with 10 machines and 50 jobs

CONCLUSION

This research build the flowshop scheduling simulator and statistics that shows the animation of the whole jobs execution from time zero until all jobs in the queue are finished to be processed. Experiment results give similar conclusions with previous paper [6] that NEH constantly perform the best compare with Palmer and Gupta algorithm. It constantly produce the minimum makespan and tardiness.

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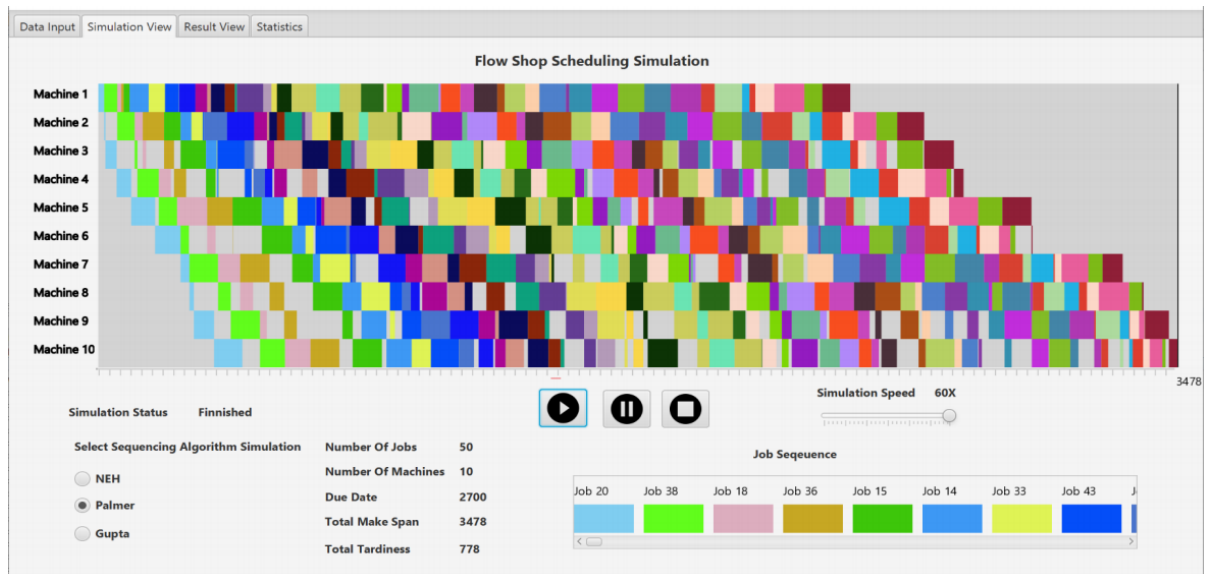


Fig. 2 Final result of Palmer execution process with 10 machines and 50 jobs

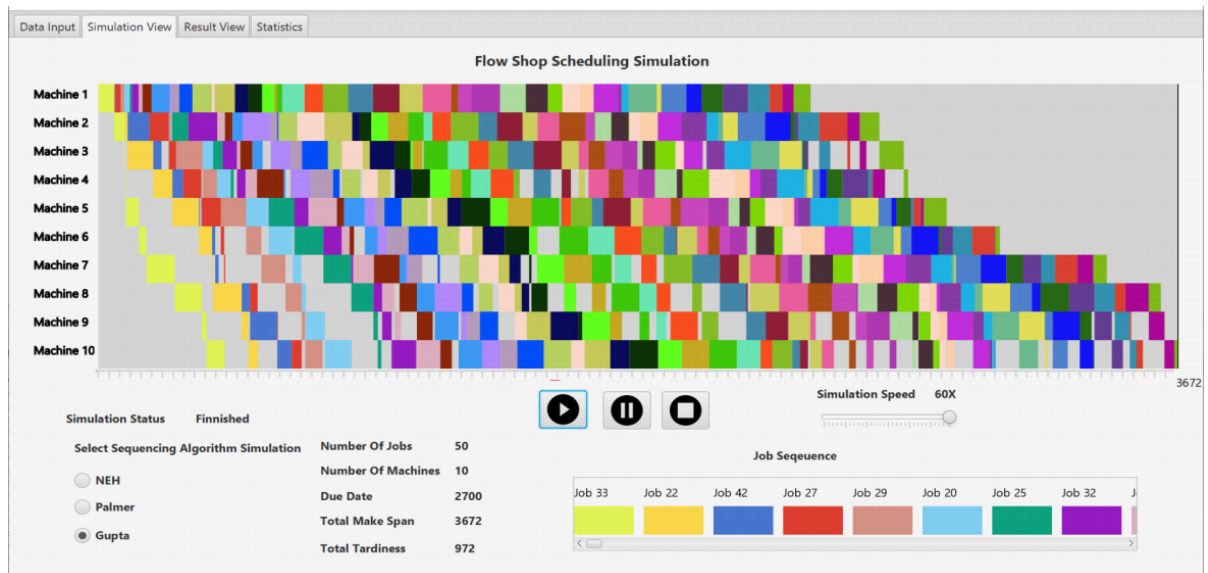


Fig. 3. Final result of Gupta execution process with 10 machines and 50 jobs

NEH Sequencing Algorithm					Palmer Sequencing Algorithm					Gupta Sequencing Algorithm				
Job	Machine 1	Machine 2	Machine 3	M	Job	Machine 1	Machine 2	Machine 3	M	Job	Machine 1	Machine 2	Machine 3	M
Job 44	38	42	23	41	Job 20	19	4	36	47	Job 33	53	37	2	2
Job 25	7	55	32	10	Job 38	42	56	9	66	Job 22	2	10	87	65
Job 30	73	15	7	54	Job 18	13	26	11	35	Job 42	2	74	28	37
Job 42	2	74	28	37	Job 36	8	69	32	35	Job 27	19	66	25	62
Job 20	19	4	36	47	Job 15	19	53	82	31	Job 29	12	7	90	45
Job 49	32	25	41	91	Job 14	61	35	35	6	Job 20	19	4	36	47
Job 32	23	99	49	25	Job 33	53	37	2	2	Job 25	7	55	32	10
Job 34	16	50	76	18	Job 43	44	4	88	22	Job 32	23	99	49	25
Job 33	53	37	2	2	Job 42	2	74	28	37	Job 18	13	26	11	35
Job 31	85	11	11	87	Job 2	52	87	1	24	Job 49	32	25	41	91
Job 36	8	69	32	35	Job 44	38	42	23	41	Job 14	61	35	35	6
Job 37	26	22	39	77	Job 29	12	7	90	45	Job 41	7	96	67	68
Job 42	2	74	28	37	Job 4	45	75	85	57	Job 37	26	22	39	77
Total Make Span 3135					Total Make Span 3478					Total Make Span 3672				

Fig. 4. Final result of NEH, Palmer, and Gupta execution process with 10 machines and 50 jobs

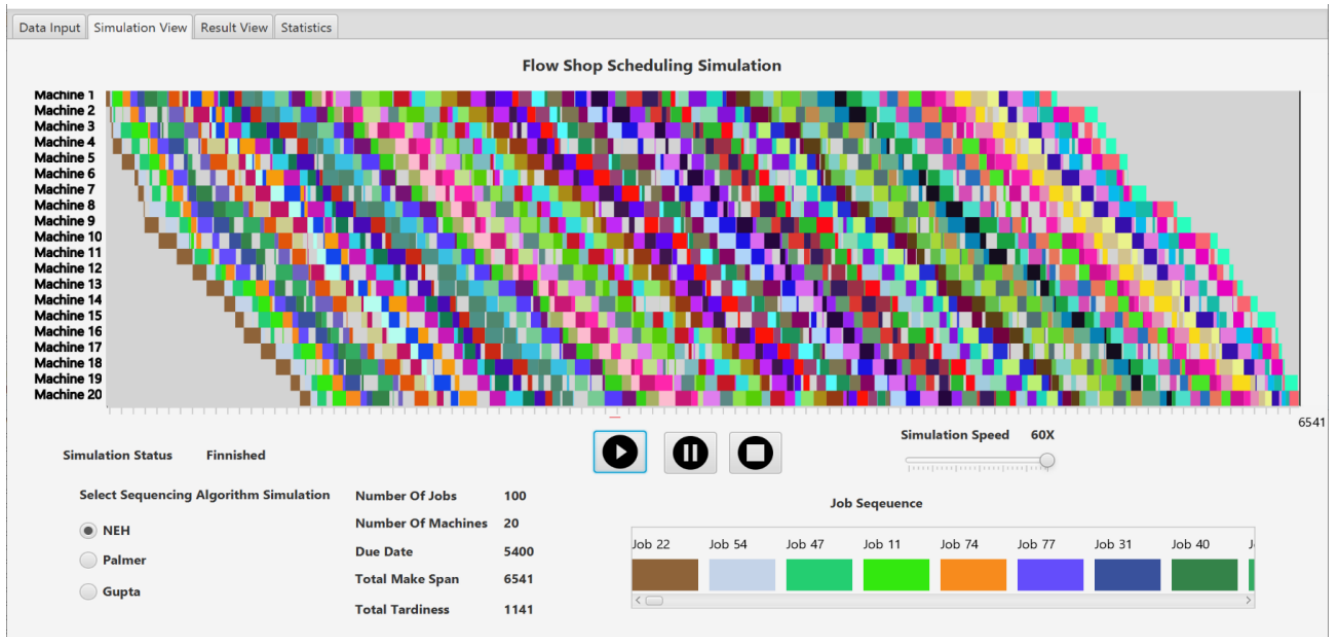


Fig. 5. Final result of Palmer execution process with 20 machines and 100 jobs



Fig. 6. Report result of NEH, Palmer, and Gupta execution process with 20 machines and 100 jobs

NEH Sequencing Algorithm					Palmer Sequencing Algorithm					Gupta Sequencing Algorithm				
Job	Machine 1	Machine 2	Machine 3	M	Job	Machine 1	Machine 2	Machine 3	M	Job	Machine 1	Machine 2	Machine 3	M
Job 22	23	8	7	45	Job 22	23	8	7	45	Job 1	12	27	24	42
Job 54	10	8	41	11	Job 33	23	16	70	83	Job 79	1	3	68	74
Job 47	5	9	59	44	Job 81	24	28	92	6	Job 41	1	86	34	73
Job 11	54	2	89	7	Job 79	1	3	68	74	Job 40	70	25	8	3
Job 74	42	46	13	40	Job 82	17	48	53	13	Job 50	33	88	10	16
Job 77	38	33	40	61	Job 54	10	8	41	11	Job 55	16	45	5	96
Job 31	40	4	16	83	Job 35	4	80	77	2	Job 98	57	12	94	23
Job 40	70	25	8	3	Job 65	62	22	48	11	Job 32	63	58	42	8
Job 24	55	3	89	20	Job 83	58	8	16	32	Job 23	47	62	99	53
Job 90	8	82	78	11	Job 9	86	42	56	6	Job 47	5	9	59	44
Job 83	58	8	16	32	Job 21	22	7	47	73	Job 22	23	8	7	45
Job 55	16	45	5	96	Job 76	77	18	7	92	Job 59	9	39	46	14
Job 50	22	88	10	16	Job 11	54	2	80	7	Job 71	2	51	82	84
Total Make Span 6541					Total Make Span 7075					Total Make Span 7741				
Total Tardiness 1141					Total Tardiness 1675					Total Tardiness 2341				

Fig. 7. Final result of Gupta execution process with 20 machines and 100 jobs

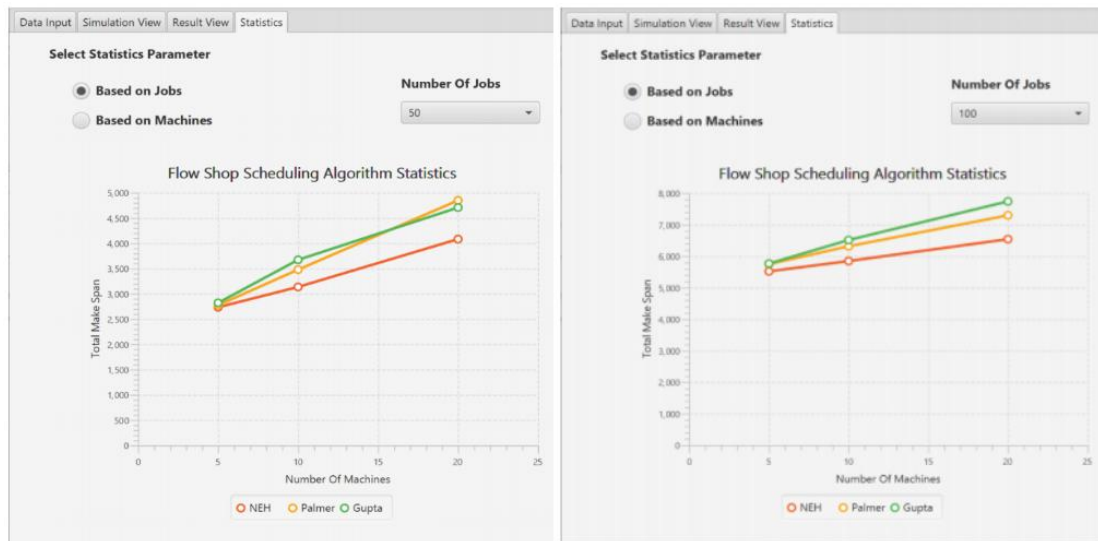


Fig. 8. Statistics of NEH, Palmer, and Gupta for 50 jobs (left) and for 100 jobs (right)